

► Stacks / Arrays / Monotonic Stack

✅ Stacks / Arrays / Monotonic Stack – Largest Rectangle in a Histogram

📌 Problem Statement:

You are given an array of bar heights where each bar has a width of `1`.

Find the **area of the largest rectangle** that can be formed inside the histogram.

📌 Example:

Input:

heights = [2, 1, 5, 6, 2, 3]

Output:

10

Explanation:

- Rectangle of height 5 and 6 → width 2 → area = $5 \times 2 = 10$

🔄 Brute Force:

- For each bar, expand left and right while height is $\geq$ current height.
- Compute area for each such range.
- **Time Complexity:** $O(n^2)$

🧠 Monotonic Stack Intuition:

- Maintain a **stack of indices** of increasing height bars.
- When a bar is **lower** than the top of the stack:
 - Pop from the stack and compute area.
 - Width = current index – index of bar in stack before the popped one – 1
- Do this until stack is empty or top is $\leq$ current bar.
- Add a dummy bar `0` at the end to flush out remaining bars.

🚀 Optimal Stack Solution

- **Time Complexity:** $O(n)$
- **Space Complexity:** $O(n)$

```
import java.util.*;
```

```

public class LargestRectangleHistogram {
    public static int largestRectangleArea(int[] heights) {
        Stack<Integer> stack = new Stack<>();
        int maxArea = 0;
        int n = heights.length;

        for (int i = 0; i <= n; i++) {
            int h = (i == n) ? 0 : heights[i];

            while (!stack.isEmpty() && h < heights[stack.peek()]) {
                int height = heights[stack.pop()];
                int width = stack.isEmpty() ? i : (i - stack.peek() - 1);
                maxArea = Math.max(maxArea, height * width);
            }

            stack.push(i);
        }

        return maxArea;
    }
}

```

✓ Output

Input: [2, 1, 5, 6, 2, 3]

Output: 10

Input: [2, 4]

Output: 4

✓ Array Manipulation / Sliding Window / Rotations – Left Rotation

🚩 Problem Statement:

Given an array of integers and a number `d`, rotate the array **to the left** by `d` positions.

Elements shifted beyond the first position move to the end.

📁 Example:

Input:

arr = [1, 2, 3, 4, 5]

d = 2

Output:

[3, 4, 5, 1, 2]



Brute Force:

- Run a loop `d` times.
 - In each iteration:
 - Remove first element.
 - Append it to the end.
 - **Time Complexity:** $O(d \times n)$
-



Better Intuition:

- Use an auxiliary array:
 - Copy elements from index `d` to `n-1`.
 - Then copy elements from index `0` to `d-1`.
-



Optimal Solution (In-place using Reversal Algorithm)

- Reverse:
 - First `d` elements.
 - Then `n-d` elements.
 - Then the whole array.
- **Time Complexity:** $O(n)$
- **Space Complexity:** $O(1)$

```
import java.util.*;

public class LeftRotation {
    public static void rotateLeft(int[] arr, int d) {
        int n = arr.length;
        d %= n;

        reverse(arr, 0, d - 1);
        reverse(arr, d, n - 1);
        reverse(arr, 0, n - 1);
    }

    private static void reverse(int[] arr, int start, int end) {
        while (start < end) {
            int temp = arr[start];
            arr[start++] = arr[end];
            arr[end--] = temp;
        }
    }
}
```

```

        arr[end--] = temp;
    }
}

public static void main(String[] args) {
    int[] arr = {1, 2, 3, 4, 5};
    int d = 2;
    rotateLeft(arr, d);
    System.out.println(Arrays.toString(arr));
}
}

```

✅ Output

Input: [1, 2, 3, 4, 5], d = 2

Output: [3, 4, 5, 1, 2]

Input: [10, 20, 30, 40], d = 1

Output: [20, 30, 40, 10]

✅ Array Manipulation – Grid Challenge

📌 Problem Statement:

You are given a square grid of characters in the form of an array of strings.

First, sort each row alphabetically.

Then, check if the **columns are also sorted** from top to bottom.

Return "YES" if columns are sorted, otherwise return "NO".

📌 Example:

Input:

```

grid = {
    "ebacd",
    "fghij",
    "olmkn",
    "trpqs",
    "xywuv"
}

```

Output:

"YES"

Explanation:

After row-wise sorting:

```
{  
    "abcde",  
    "fghij",  
    "klmno",  
    "pqrst",  
    "uvwxy"  
}
```

All columns are sorted.



Brute Force:

- Sort each row.
- For each column, check if it's sorted.
- **Time Complexity:** $O(n^2 \log n)$ for sorting rows + $O(n^2)$ for column check.



Optimal Solution (Straightforward)

- Sort each row.
- For each column `c`, check if `grid[i][c] <= grid[i+1][c]`
- **Time Complexity:** $O(n^2 \log n)$



Java Code:

```
public class GridChallenge {  
    public static String gridChallenge(String[] grid) {  
        int n = grid.length;  
  
        // Sort each row  
        for (int i = 0; i < n; i++) {  
            char[] chars = grid[i].toCharArray();  
            Arrays.sort(chars);  
            grid[i] = new String(chars);  
        }  
  
        // Check columns  
        for (int col = 0; col < grid[0].length(); col++) {  
            for (int row = 0; row < n - 1; row++) {  
                if (grid[row].charAt(col) > grid[row + 1].charAt(col)) {  
                    return "NO";  
                }  
            }  
        }  
    }  
}
```

```

        return "YES";
    }

    public static void main(String[] args) {
        String[] grid = {"ebacd", "fghij", "olmkn", "trpqs", "xywuv"};
        System.out.println(gridChallenge(grid)); // Output: YES
    }
}

```

✓ Output

Input:

`["ebacd", "fghij", "olmkn", "trpqs", "xywuv"]`

Output: `"YES"`

Input:

`["zyx", "wvu", "tsr"]`

Output: `"NO"`

✓ Array Manipulation / Greedy – Mark and Toys

🚧 Problem Statement:

Mark wants to buy as many toys as possible with a given budget.

You are given a list of toy prices and an integer budget `k`.

Determine the **maximum number of toys** he can purchase.

📖 Example:

Input:

`prices = [1, 12, 5, 111, 200, 1000, 10]`

`k = 50`

Output:

`4`

Explanation:

He can buy `[1, 5, 10, 12]` = 28 total spent → max 4 toys.

🔄 Brute Force:

- Generate all subsets and check which ones fit under budget.
- **Time Complexity:** Exponential → Not practical.

Greedy Intuition:

- Sort the prices.
 - Start buying the **cheapest toys first** until the budget runs out.
-

Optimal Greedy Solution

- Sort prices.
 - Keep a running sum.
 - **Time Complexity:** $O(n \log n)$
-

Java Code:

```
import java.util.*;

public class MarkAndToys {
    public static int maximumToys(int[] prices, int k) {
        Arrays.sort(prices);
        int count = 0;
        int total = 0;

        for (int price : prices) {
            if (total + price <= k) {
                total += price;
                count++;
            } else {
                break;
            }
        }

        return count;
    }

    public static void main(String[] args) {
        int[] prices = {1, 12, 5, 111, 200, 1000, 10};
        int k = 50;
        System.out.println(maximumToys(prices, k)); // Output: 4
    }
}
```

Output

Input:

prices = [1, 12, 5, 111, 200, 1000, 10], k = 50

Output: 4

✓ Array Manipulation / Greedy – Permuting Two Arrays

📌 Problem Statement:

You are given two arrays **A** and **B**, both of length **n**, and an integer **k**.

Can you permute both arrays such that **$A[i] + B[i] \geq k$** for all **i**?

Return **"YES"** if it's possible, otherwise return **"NO"**.

📖 Example:

Input:

A = [2, 1, 3]

B = [7, 8, 9]

k = 10

Output:

"YES"

Explanation:

Sorted A = [1, 2, 3], B = [9, 8, 7]

$A[0] + B[0] = 10$, $A[1] + B[1] = 10$, $A[2] + B[2] = 10$ ✓

🔄 Brute Force:

- Try all permutations of both arrays.
- For each permutation, check the sum condition.
- **Time Complexity:** $O(n!^2)$ — not practical.

🚀 Optimal Greedy Solution:

- Sort A in ascending order.
- Sort B in **descending** order.
- Check if $A[i] + B[i] \geq k$ for all **i**.
- **Time Complexity:** $O(n \log n)$

✓ Java Code:

```
import java.util.*;

public class PermutingTwoArrays {
    public static String twoArrays(int k, int[] A, int[] B) {
```



```

Arrays.sort(A);
Arrays.sort(B);
int n = A.length;

// Reverse B to make it descending
for (int i = 0; i < n / 2; i++) {
    int temp = B[i];
    B[i] = B[n - 1 - i];
    B[n - 1 - i] = temp;
}

for (int i = 0; i < n; i++) {
    if (A[i] + B[i] < k) return "NO";
}
return "YES";
}

public static void main(String[] args) {
    int[] A = {2, 1, 3};
    int[] B = {7, 8, 9};
    int k = 10;
    System.out.println(twoArrays(k, A, B)); // Output: YES
}
}

```

✓ Output

Input:

A = [2, 1, 3], B = [7, 8, 9], k = 10

Output:

"YES"

✓ Heap / Priority Queue – Jesse and Cookies

🔥 Problem Statement:

Jesse loves cookies. Each cookie has a **sweetness level**. Jesse wants to **combine** the two least sweet cookies until all cookies have **at least sweetness** `k`.

To combine two cookies with sweetness `a` and `b`:

New cookie = `a + 2*b`

Determine the **minimum number of operations** required.

Return `-1` if it's not possible.

Example:

Input:

cookies = [1, 2, 3, 9, 10, 12]

k = 7

Output:

2

Explanation:

- Combine 1 + 2×2 = 5 → [3, 5, 9, 10, 12]

- Combine 3 + 2×5 = 13 → [9, 10, 12, 13]

Now all ≥ 7 

Brute Force:

- Sort array on every iteration.
- Pick smallest 2, combine, insert again.
- **Time Complexity:** $O(n^2 \log n)$

Optimal (Min Heap):

- Use a **min-heap** (priority queue).
- Pop 2 smallest elements.
- Push combined cookie.
- Repeat until heap top ≥ k or heap size < 2.
- **Time Complexity:** $O(n \log n)$

Java Code:

```
import java.util.*;

public class JesseAndCookies {
    public static int cookies(int k, int[] A) {
        PriorityQueue<Integer> minHeap = new PriorityQueue<>();
        for (int num : A) minHeap.add(num);

        int operations = 0;

        while (minHeap.size() > 1 && minHeap.peek() < k) {
            int least = minHeap.poll();
            int secondLeast = minHeap.poll();
            int combined = least + 2 * secondLeast;
```

```

        minHeap.add(combined);
        operations++;
    }

    return (minHeap.peek() >= k) ? operations : -1;
}

public static void main(String[] args) {
    int[] cookies = {1, 2, 3, 9, 10, 12};
    int k = 7;
    System.out.println(cookies(k, cookies)); // Output: 2
}
}

```

✅ Output

Input:

cookies = [1, 2, 3, 9, 10, 12], k = 7

Output: 2

✅ Heap / Priority Queue – Find the Running Median

🔗 Problem Statement:

Given a stream of integers, for each new element added, print the **median** of all elements seen so far.

📌 Example:

Input stream: [12, 4, 5, 3, 8, 7]

Output:

12

8

5

4.5

5

6

😬 Intuition:

To efficiently compute the **running median**:

- Use **two heaps**:
 - Max-heap for **lower half**
 - Min-heap for **upper half**

- Ensure:
 - Size difference ≤ 1
 - Max-heap top $\leq$ Min-heap top

Median is:

- If equal size $\rightarrow$ average of tops
- Else $\rightarrow$ top of the bigger heap

Brute Force:

- Sort the entire array on each insertion.
- Find the middle element.
- **Time Complexity:** $O(n^2 \log n)$

Optimal Heap-Based Solution:

- **MaxHeap:** for smaller half
- **MinHeap:** for larger half
- Keep heaps balanced.
- **Time Complexity:** $O(n \log n)$

Java Code:

```
import java.util.*;

public class RunningMedian {
    public static List<Double> runningMedian(int[] nums) {
        List<Double> result = new ArrayList<>();
        PriorityQueue<Integer> maxHeap = new PriorityQueue<>
(Collections.reverseOrder()); // left
        PriorityQueue<Integer> minHeap = new PriorityQueue<>(); // right

        for (int num : nums) {
            if (maxHeap.isEmpty() || num <= maxHeap.peek()) {
                maxHeap.offer(num);
            } else {
                minHeap.offer(num);
            }

            // Balance heaps
            if (maxHeap.size() > minHeap.size() + 1) {
                minHeap.offer(maxHeap.poll());
            } else if (minHeap.size() > maxHeap.size()) {
```

```

        maxHeap.offer(minHeap.poll());
    }

    // Get median
    if (maxHeap.size() == minHeap.size()) {
        result.add((maxHeap.peek() + minHeap.peek()) / 2.0);
    } else {
        result.add((double) maxHeap.peek());
    }
}

return result;
}

public static void main(String[] args) {
    int[] nums = {12, 4, 5, 3, 8, 7};
    List<Double> medians = runningMedian(nums);
    for (Double m : medians) {
        System.out.println(m);
    }
}
}

```

✓ Output:

```

12.0
8.0
5.0
4.5
5.0
6.0

```